Exercises EMGMT: Algorithmic Analysis

Note: All logarithms have base 2.

- 1. What is the sum of all even numbers from 0 to 2n, for any integer n > 0?
- 2. Count the number of primitive operations that are performed in each of the two algorithms below.

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    (a) Algorithm arrayMax(A,n):
    Input: An array A storing n ≥ 1 integers.
    Output: The maximum element in A.
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\begin{array}{l} \operatorname{currentMax} \leftarrow \operatorname{A}[0] \\ \operatorname{\textbf{for}} i \leftarrow 1 \text{ to } n-1 \text{ do} \\ \operatorname{\textbf{for}} j \leftarrow i \text{ to } n-1 \text{ do} \\ \operatorname{\textbf{if}} \operatorname{A}[i] > \operatorname{A}[j] \text{ then} \\ \operatorname{\textbf{if}} \operatorname{A}[i] > \operatorname{currentMax} \text{ then} \\ \operatorname{currentMax} \leftarrow \operatorname{A}[i] \end{array}
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 $\mathbf{return} \ \mathbf{current} \mathbf{Max}$

(b) Algorithm weird(A,n): Input: An array A storing $n \ge 1$ integers.

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Output: An array F storing n \ge 1 integers.
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for i \leftarrow 1 to n do

F[i] \leftarrow 0

j \leftarrow i

while j \le n do

for k \leftarrow 0 to 4 do

F[i] \leftarrow F[i] + A[j + k]

j \leftarrow 2 \cdot j

return F
```

- 3. For each of the items below, find the smallest $n_0 \in \mathbb{N}$ such that for all $n \geq n_0$:
 - (a) n^3 is greater than $4n^2 + 60n$,
 - (b) $8n \log(n)$ is smaller than $2n^2$,
 - (c) 2^n is greater than n^4 .
- 4. Give a c > 0 and an integer $n_0 \ge 1$ such that, for all $n \ge n_0$:
 - (a) $40n^2 + 120 \le cn^3$,
 - (b) $16n \log(n^2) \le cn^2$,
 - (c) $\frac{1}{10}2^n \ge cn^4$.

- 5. Give a function f(n) such that $22n^4 \log(n) + 130 \log^2(n)$ is $\Theta(f(n))$.
- 6. Order the following functions by asymptotic growth rate:

 $\begin{array}{rrrr} 4n\log(n)+2n & 2^{10} & 2^{\log(n)} \\ 3n+100\log(n) & 4n & 2^n \\ n^2+10n & n^3 & n\log(n) \end{array}$

7. Prove:

(a) $2n^3 + 9n^2$ is $O(n^3)$

(b) $(n \log n)/8$ is $\Omega(n \log n)$

(c) $2^{n+2} - n$ is $\Theta(2^n)$

- 8. Prove or disprove: if d(n) is O(f(n)), then ad(n) is O(f(n)), for any constant a > 0.
- 9. Prove or disprove: if d(n) is O(f(n)) and e(n) is O(g(n)), then d(n) e(n) is O(f(n) g(n)).
- 10. What are the running times of the algorithms in exercises 2(a) and 2(b) in big-Oh-notation?